

# Quantum Computing: From Circuit To Architecture

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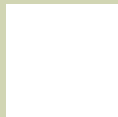
This talk is mainly based on portions of the course held by Carmen G. Almudever (TU Delft) at Acaces summer school (Fiuggi, 9-14th July 2017), enriched with some material by myself

## Outline

1. Qu-bit definition
2. Quantum gates
3. Multi-states
4. Example: Teleportation
5. Quantum Algorithms
6. Quantum Processor
7. Compilation
8. Quantum Computers Architecture

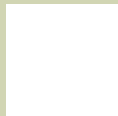
## Classical Bit

Bit is either 0 or 1:



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Bit is either 0 or 1:



## Quantum Bit

The bit is in a superposition state: it is both 0 and 1



A qu-bit  $\psi$  is defined as:  $|\psi\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle$       $\alpha_0, \alpha_1 \in \mathbb{C}$

$$\text{Since } |0\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \text{ and } |1\rangle = \begin{bmatrix} 0 \\ 1 \end{bmatrix} \Rightarrow |\psi\rangle = \begin{bmatrix} \alpha_0 \\ \alpha_1 \end{bmatrix}$$

# Qu-Bit: Measurement

What do  $\alpha_0$  and  $\alpha_1$  actually mean?

## Measurement

Consider a qu-bit  $|\psi\rangle = [\alpha_0 \quad \alpha_1]$ . Define the measurement as a function  $M(|\psi\rangle)$  with range  $\{0, 1\}$ , such that:

- $Pr(M(|\psi\rangle) = 0) = |\alpha_0|^2$
- $Pr(M(|\psi\rangle) = 1) = |\alpha_1|^2$

Therefore, it must be  $|\alpha_0|^2 + |\alpha_1|^2 = 1$



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The superposition state is destroyed after measurement!



We cannot directly measure the superposition, only probabilistic estimation

# Qu-Bit: A Real World Example





## Crystal of Tourmaline: Classical World

- Interaction with plane-polarized light:
  1. Light polarized perpendicularly w.r.t. the crystal axis  $\Rightarrow$  The light goes through the crystal
  2. Light polarized parallel w.r.t. the crystal axis  $\Rightarrow$  The light is filtered by the crystal
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## Crystal of Tourmaline: Quantum World

- Interaction with a single plane-polarized photon:
  1. Photon polarized perpendicularly w.r.t. the crystal axis  $\Rightarrow$  The photon is detected after the crystal
  2. Photon polarized parallel w.r.t. the crystal axis  $\Rightarrow$  The photon is not detected after the crystal
  3. Photon polarized with angle  $\alpha$  w.r.t. the crystal axis  $\Rightarrow$  A photon perpendicularly polarized is detected  $\sin^2 \alpha$  times, no photon detected otherwise

## From Physic World to Qu-Bit

Qu-bit  $\Leftrightarrow$  the polarization direction of a single photon

$|0\rangle$  photon polarized perpendicular w.r.t. the crystal axis

$|1\rangle$  photon polarized parallel w.r.t. the crystal axis

Superposition state? a photon polarized with angle  $\alpha$  w.r.t. the crystal axis:

$$|\psi\rangle = \sin \alpha |0\rangle + \cos \alpha |1\rangle$$

## Measurement



- The qu-bit is 0 with probability  $\sin^2 \alpha$
- The qu-bit is 1 with probability  $\cos^2 \alpha$
- The qu-bit is destroyed: no longer polarized with angle  $\alpha$

- Qu-bits are vectors in  $\mathbb{C} \Rightarrow$  Gates are matrices in  $\mathbb{C}$
- Properties? Unitary Operations!
- Generic gate  $\text{---} \boxed{U} \text{---} : UU^* = U^*U = I \Rightarrow |\det(U)| = 1$



Quantum gates are reversible!

# Quantum Gates

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## Main Single Qu-Bit Gates

- Bit Flip Gate:  $\text{---} \boxed{X} \text{---} \Rightarrow \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$
- Identity Gate:  $\text{---} \boxed{I} \text{---} \Rightarrow \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$
- Phase Flip Gate:  $\text{---} \boxed{Z} \text{---} \Rightarrow \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$
- Hadamar Gate:  $\text{---} \boxed{H} \text{---} \Rightarrow \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$

## Hadamar Gate Effect

$$|\psi_{out}\rangle = H |\psi_{in}\rangle, \quad H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$|\psi_{in}\rangle = |0\rangle$$

$$\begin{aligned} |\psi_{out}\rangle &= H|0\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \\ &= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \end{aligned}$$

$$|\psi_{in}\rangle = |1\rangle$$

$$\begin{aligned} |\psi_{out}\rangle &= H|1\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \\ &= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 \\ -1 \end{bmatrix} = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \end{aligned}$$

$\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) = |+\rangle$  and  $\frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) = |-\rangle$  are 2 relevant states.  
Why?

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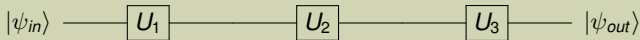
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Why?



For both of them,  $|\alpha_0|^2 = |\alpha_1|^2 = \frac{1}{2}$

## From Gates to Circuits



$$|\psi_{out}\rangle = U_3 U_2 U_1 |\psi_{in}\rangle$$

Reversibility: The way back!



$$|\psi_{in}\rangle = U_1^* U_2^* U_3^* |\psi_{out}\rangle$$



## Examples

$$\begin{aligned}
 |0\rangle \text{ --- } \boxed{H} \text{ --- } \boxed{Z} \text{ --- } \boxed{H} &= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \\
 &= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & -1 \\ 1 & 1 \end{bmatrix} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \\
 &= \frac{1}{2} \begin{bmatrix} 0 & 2 \\ 2 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = X|0\rangle \\
 &= |1\rangle
 \end{aligned}$$

$$\begin{aligned}
 |1\rangle \text{ --- } \boxed{H} \text{ --- } \boxed{X} \text{ --- } \boxed{H} &= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \\
 &= \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ -1 & 1 \end{bmatrix} \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} \\
 &= \frac{1}{2} \begin{bmatrix} 2 & 0 \\ 0 & -2 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = Z|1\rangle \\
 &= -|1\rangle
 \end{aligned}$$

- Concise representation of superposition of multiple bits
- Real computational power of quantum computers!
- We need to introduce a new algebraic operation: the tensor product  $\otimes$

# Multi Qu-Bit State

- Concise representation of superposition of multiple bits
- Real computational power of quantum computers!
- We need to introduce a new algebraic operation: the tensor product  $\otimes$

## Tensor Product

Given  $A = \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}$  and  $B \in \mathbb{C}^{n \times m}$ , the tensor product is defined as:

$$T = A \otimes B = \begin{bmatrix} a_{11}B & a_{12}B \\ a_{21}B & a_{22}B \end{bmatrix}$$

Example:  $A = \begin{bmatrix} 2 \\ 3 \end{bmatrix}$        $B = \begin{bmatrix} 1 & 4 \\ 5 & 2 \end{bmatrix}$

$$T = \begin{bmatrix} 2 \begin{bmatrix} 1 & 4 \\ 5 & 2 \end{bmatrix} \\ 3 \begin{bmatrix} 1 & 4 \\ 5 & 2 \end{bmatrix} \end{bmatrix} = \begin{bmatrix} 2 & 8 \\ 10 & 4 \\ 3 & 12 \\ 15 & 6 \end{bmatrix}$$

- The superposition state may apply to multiple qu bits
- Instead of the 2 kets  $|0\rangle$  and  $|1\rangle$ , there is a ket for each possible combination of bits
- The coefficients are related to the measurement probability of the corresponding combination

## 2 Qu-bits State

$$|\psi\rangle = \alpha_0 |00\rangle + \alpha_1 |01\rangle + \alpha_2 |10\rangle + \alpha_3 |11\rangle$$
$$|\alpha_0|^2 + |\alpha_1|^2 + |\alpha_2|^2 + |\alpha_3|^2 = 1$$

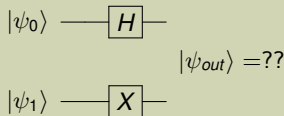
Vector representation?  $\Rightarrow$  tensor product between single qu-bit kets

$$|00\rangle = |0\rangle \otimes |0\rangle = \begin{bmatrix} 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 0 \end{bmatrix} = \begin{bmatrix} 1 & [1 & 0] & 0 & [1 & 0] \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \end{bmatrix}$$

$$|10\rangle = |1\rangle \otimes |0\rangle = \begin{bmatrix} 0 & 1 \end{bmatrix} \otimes \begin{bmatrix} 1 & 0 \end{bmatrix} = \begin{bmatrix} 0 & [1 & 0] & 1 & [1 & 0] \end{bmatrix} = \begin{bmatrix} 0 & 0 & 1 & 0 \end{bmatrix}$$

$$|\psi_2\rangle = |\psi_0\rangle \otimes |\psi_1\rangle = \begin{bmatrix} \alpha_0 & \alpha_1 \end{bmatrix} \otimes \begin{bmatrix} \beta_0 & \beta_1 \end{bmatrix} = \begin{bmatrix} \alpha_0 & [\beta_0 & \beta_1] & \alpha_1 & [\beta_0 & \beta_1] \end{bmatrix}$$
$$= \begin{bmatrix} \alpha_0\beta_0 & \alpha_0\beta_1 & \alpha_1\beta_0 & \alpha_1\beta_1 \end{bmatrix}$$

## 2 Qu-bits Circuit



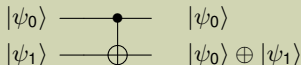
- The gates operate on the multi qu-bit  $|\psi_{in}\rangle = |\psi_0\rangle \otimes |\psi_1\rangle$
- To apply the gates, we need to combine them in a 2 qu-bit gate. How?
- Tensor Product!

$$|\psi_{out}\rangle = (H \otimes X) |\psi_{in}\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} & \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \\ \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} & \begin{bmatrix} 0 & -1 \\ -1 & 0 \end{bmatrix} \end{bmatrix}$$

$$|\psi_{in}\rangle = |00\rangle \rightarrow |\psi_{out}\rangle = \frac{1}{\sqrt{2}} \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & -1 \\ 1 & 0 & -1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ \frac{1}{\sqrt{2}} \\ 0 \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \frac{1}{\sqrt{2}}(|01\rangle + |11\rangle)$$

There are also quantum gates which apply only on multiple qu-bits

## CNOT Gate



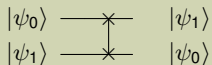
- If the control bit ( $|\psi_0\rangle$ ) is 1, then the target bit ( $|\psi_1\rangle$ ) is inverted

- The gate matrix is already  $4 \times 4$ :
$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

- Superposition state: act independently on the fundamental states

↪ Example:  $|\psi_{in}\rangle = \frac{\sqrt{3}}{2} |00\rangle + \frac{1}{2} |10\rangle \rightarrow |\psi_{out}\rangle = \frac{\sqrt{3}}{2} |00\rangle + \frac{1}{2} |11\rangle$

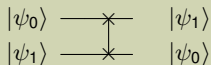
## Swap Gate



- The Qu-Bits are swapped

- Gate Matrix: 
$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

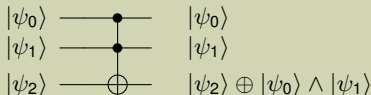
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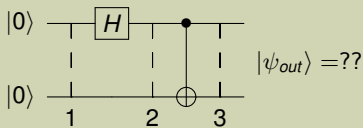
## Toffoli Gate



- CNOT gate with 2 control bits instead of 1
- That is, the target bit ( $|\psi_2\rangle$ ) is inverted when both control bits are 1



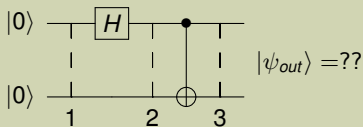
## Example



Step-by-Step evaluation:

1.  $|\psi_{in}\rangle = |00\rangle$
2.  $\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)|0\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |10\rangle)$
3.  $|\psi_{out}\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$

## Example

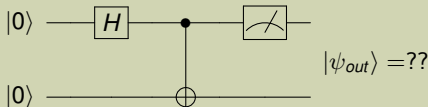


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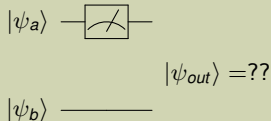
## Measuring 1 Qu-Bit

Introduce a slight variation of the above circuit:



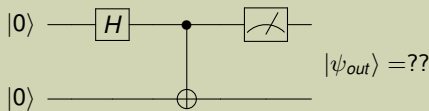
What happens to the multi qu-bits state when not all the bits are measured?

## Partial Measurement: 2 Qu-Bits State



- Before measurement, the multi qu-bit  $|\psi_{in}\rangle = |\psi_a\rangle \otimes |\psi_b\rangle$
- As with single bit gates, bit-wise reasoning
  1. Splitting qu-bits: Given a 2 multi qu-bit state  $|\psi\rangle$ , it is always possible to split it as  $|\psi\rangle = \alpha_0 |0\rangle \otimes (|\psi_0\rangle) + \alpha_1 |1\rangle \otimes (|\psi_1\rangle)$ , such that  $|\alpha_0|^2 + |\alpha_1|^2 = 1$  and  $|\psi_0\rangle, |\psi_1\rangle$  are valid qu-bits
  2. Then, depending on the measurement outcome on the first qu-bit:
    - Measurement of the first qu-bit is 0  $\rightarrow$  the second qu-bit is  $|\psi_0\rangle \rightarrow |\psi_{out}\rangle = |\psi_0\rangle$
    - Measurement of the first qu-bit is 1  $\rightarrow$  the second qu-bit is  $|\psi_1\rangle \rightarrow |\psi_{out}\rangle = |\psi_1\rangle$

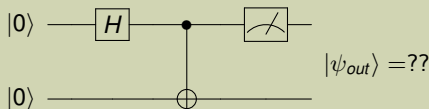
## Back to the Example Circuit



Before measurement, the multi qu-bit  $|\psi\rangle = \frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$

1.  $|\psi\rangle = \frac{1}{\sqrt{2}} |0\rangle \otimes |0\rangle + \frac{1}{\sqrt{2}} |1\rangle \otimes |1\rangle$
2.  $\psi_{out} = |0\rangle$  if the measured qu-bit is 0
3.  $\psi_{out} = |1\rangle$  if the measured qu-bit is 1

## Back to the Example Circuit

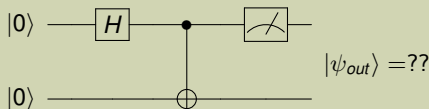


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- The measurement of a bit determines the second qu-bit
- This is weird, since the measurement affects only 1 bit

## Back to the Example Circuit



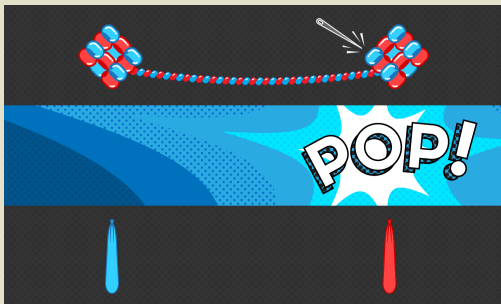
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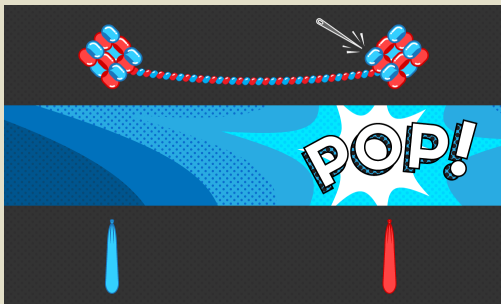
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- The measurement of a bit determines the second qu-bit
- This is weird, since the measurement affects only 1 bit



A quantum phenomenon is happening: **entanglement!**





## Entanglement Definition

- Recall the splitting of a 2 multi qu-bit state  $|\psi\rangle = |\psi_a\rangle \otimes |\psi_b\rangle$ :  
 $|\psi\rangle = \alpha_0 |0\rangle \otimes (|\psi_0\rangle) + \alpha_1 |1\rangle \otimes (|\psi_1\rangle)$ , such that  $|\alpha_0|^2 + |\alpha_1|^2 = 1$  and  $|\psi_0\rangle, |\psi_1\rangle$  are valid qu-bits
- $|\psi_a\rangle$  and  $|\psi_b\rangle$  are entangled  $\Leftrightarrow |\psi_0\rangle \neq |\psi_1\rangle$
- Meaning: the quantum state of the unmeasured bit depends on the measurement outcome of the entangled bit.



# Entanglement: Examples

- $$|\psi\rangle = \frac{1}{2}(|00\rangle + |01\rangle + |10\rangle + |11\rangle) =$$

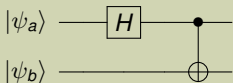
$$\frac{1}{\sqrt{2}}|0\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) + \frac{1}{\sqrt{2}}|1\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \rightarrow \text{Not entangled!}$$
- $$|\psi\rangle = \frac{1}{2}(|00\rangle + |01\rangle + |10\rangle - |11\rangle) =$$

$$\frac{1}{\sqrt{2}}|0\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) + \frac{1}{\sqrt{2}}|1\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \rightarrow \text{Entangled!}$$

# Entanglement: Examples

- $|\psi\rangle = \frac{1}{2}(|00\rangle + |01\rangle + |10\rangle + |11\rangle) =$   
 $\frac{1}{\sqrt{2}}|0\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) + \frac{1}{\sqrt{2}}|1\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) \rightarrow$  Not entangled!
- $|\psi\rangle = \frac{1}{2}(|00\rangle + |01\rangle + |10\rangle - |11\rangle) =$   
 $\frac{1}{\sqrt{2}}|0\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle) + \frac{1}{\sqrt{2}}|1\rangle \otimes \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \rightarrow$  Entangled!

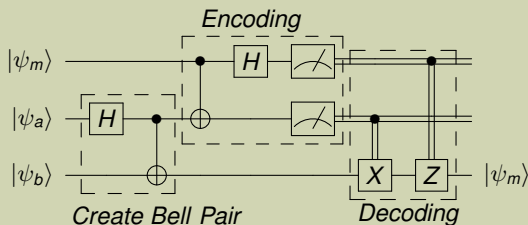
## Bell States



$|\psi_a\rangle \otimes |\psi_b\rangle = |\psi_{in}\rangle \in \{|00\rangle, |01\rangle, |10\rangle, |11\rangle\} \Rightarrow |\psi_a\rangle$  and  $|\psi_b\rangle$  are entangled at the end of the circuit:

$ \psi_a\rangle$	$ \psi_b\rangle$	
	$ 0\rangle$	$ 1\rangle$
$ 0\rangle$	$\frac{1}{\sqrt{2}}( 00\rangle +  11\rangle)$	$\frac{1}{\sqrt{2}}( 01\rangle +  10\rangle)$
$ 1\rangle$	$\frac{1}{\sqrt{2}}( 00\rangle -  11\rangle)$	$\frac{1}{\sqrt{2}}( 01\rangle -  10\rangle)$

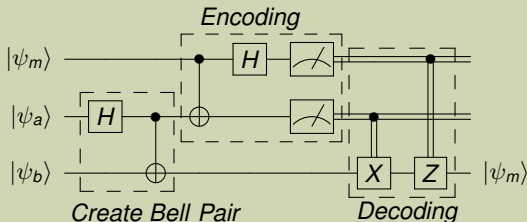
## Teleportation Circuit



- The sender and the receiver generates a bell pair
- The sender keeps  $|\psi_a\rangle$ , while the receiver keeps  $|\psi_b\rangle$
- When the sender wants to send a qu-bit  $|\psi_m\rangle$ , it performs encoding using  $|\psi_a\rangle$  too
- 2 classical bits are sent to the receiver for the decoding procedure
- After decoding, the entangled qu-bit  $|\psi_b\rangle$  has become equal to  $|\psi_m\rangle$

# Quantum Teleportation

## Teleportation Circuit



Example with  $|\psi_m\rangle = |1\rangle$  and  $|\psi_a\rangle = |\psi_b\rangle = |0\rangle \rightarrow |\psi_{in}\rangle = |100\rangle$ :

1. Bell pair creation:  $\frac{1}{\sqrt{2}}(|100\rangle + |111\rangle)$
2. Encoding before measurement:  $\frac{1}{2}(|010\rangle - |110\rangle + |001\rangle - |101\rangle) = \frac{1}{2}(|00\rangle \otimes |1\rangle + |01\rangle \otimes |0\rangle + |10\rangle \otimes (-|1\rangle) + |11\rangle \otimes (-|0\rangle))$

3. Decoding:

Measure	$ \psi_b\rangle$	Corrections	Output
00	$ 1\rangle$	No	$ 1\rangle$
01	$ 0\rangle$	X	$ 1\rangle$
10	$- 1\rangle$	Z	$ 1\rangle$
11	$- 0\rangle$	X,Z	$ 1\rangle$

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## Qu-Bit Power

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- $N$  classical bits hold 1 single value between 0 and  $2^N - 1$
- For instance, 010 is the value 2
- Classical computation performs only on the single value of the bits

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### Quantum Bits:

- $N$  qu-bits contains all possible  $2^N$  values representable by  $N$  bits
- For instance,  
 $\alpha_0 |000\rangle + \alpha_1 |001\rangle + \alpha_2 |010\rangle + \alpha_3 |011\rangle + \alpha_4 |100\rangle + \alpha_5 |101\rangle + \alpha_6 |110\rangle + \alpha_7 |111\rangle$  represents all integers from 0 to 7
- Performing quantum computation is equivalent to compute at the same time with all these  $2^N$  values. How?  $\Rightarrow$  **Quantum Algorithms!**

## Description

Quantum algorithms structure:

- Work on multi qu-bits in superposition states
- The operations performed on the qu-bits are chosen to get to a final superposition state
- Measuring this state generally yields the solution of the problem with probability close to 1
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## Example: Integer Factorization

Classical computing for a 2048 bits number?



- 100 years
- $10^5$  Trillion €
- 398549  $km^2$  server farm

Quantum computation?  $\Rightarrow$  26.7 hours using Shor algorithm!

## Description

- Classical reduction to the order-finding subproblem
- This sub-problem is solved with the quantum algorithm
- Some properties of the solution are tested, otherwise the procedure is repeated to yield a new solution of the subproblem

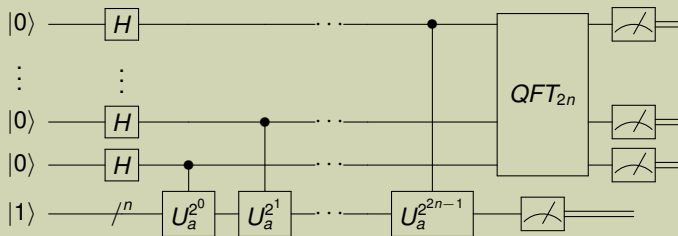
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## Order-Finding Solver

Given  $f(x) = a^x \pmod N$ , find the period of  $f$ , i.e. the order of  $a$



Quantum computing is extremely powerful, but . . .



Quantum Technologies are extremely fragile!

Quantum computing is extremely powerful, but . . .



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## Quantum Computation Errors

- A qu-bit is affected by external noise
- For instance, the Brownian motions of the molecules may interfere with the quantum estate
- Each qu-bit has a decoherence time: The maximum time a qu-bit can keep its superposition state
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Quantum technology fragility



Error correction codes are necessary to preserve the computation

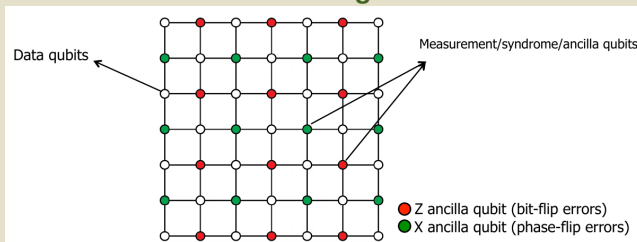
## Quantum Error Correction Codes

- Correction process is carried on after each operation
- As every correction code, redundancy is employed to correct errors
- A lot of redundancy is necessary, since:
  1. Qu-bits are continuous, not discrete
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- Each qu-bit becomes a **logical** qu-bit, which is encoded in  $n$  **physical** qu-bits: **data** and **ancilla** ones

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### Surface Code Logical Qu-Bit:





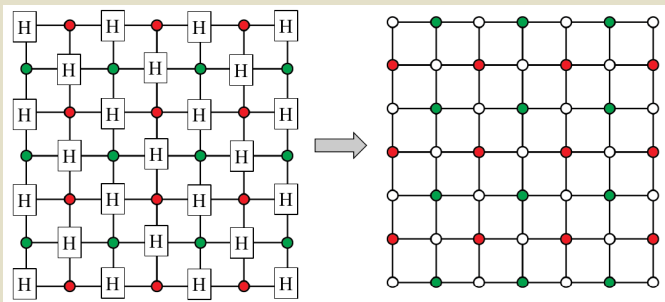
## Quantum Operations on Logical Qu-Bits

- Each of the fundamental operations ( $X, Z, H, \dots$ ) needs to be defined on the logical qu-bit
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### Surface Code Hadamar Gate:



Error correction is the main responsible for the blowup of qu-bits required for a quantum algorithm

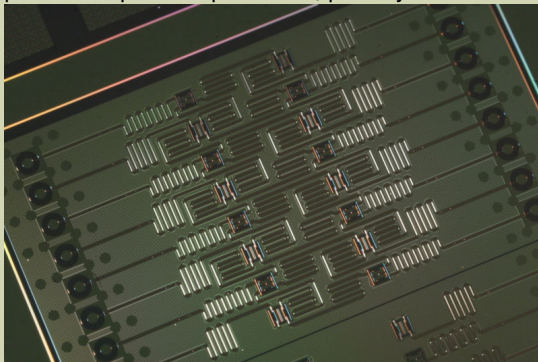
## Shor Algorithm Overhead

For instance, Shor algorithm on  $L = 2048$  bits number requires:

Rationale	#Physical Qu-bits (cumulative)
$6L$ logical qu-bits	12,288
$8 \times$ ancilla qu-bits	98,304
$1.33 \times$ to provide 'wiring' room to move qu-bits	133,000
$10k \times$ surface code	$1.3bn$
4x micro-architecture details	$5.2bn$

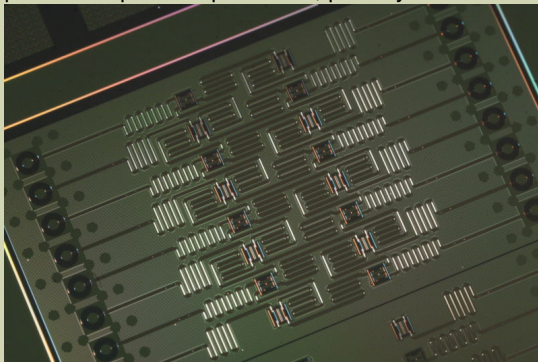
## How Many Qu-Bits?

16 qu-bit IBM quantum processor, publicly available online:



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IBM Trend:

TECH—

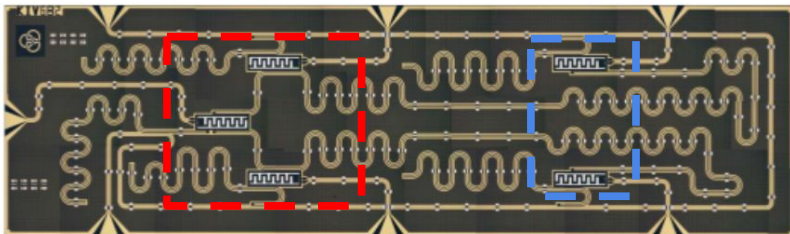
**IBM will sell 50-qubit universal quantum computer “in the next few years”**

IBM has solved most of the science behind quantum computing. Time to make some money.

SEBASTIAN ANTHONY - 6/3/2017, 12:59

## Quantum Chips

5 qu-bits chip scheme for a logical qu-bit:



- data qu-bits, ancilla qu-bits
- The gates are implemented via microwave pulses ( $10^{-8}$ s) sent to the qu-bits
- We want to perform different gates within the decoherence time

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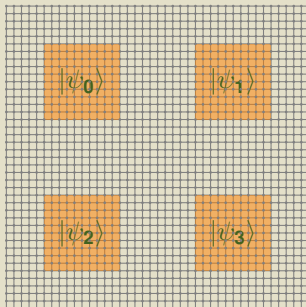


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- Optimization** Use heuristics to merge gates or re-arrange operations
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- Quantum Execution:** Translation of gate-level instructions (QISA) to signals sent to the processor

Logical view of the chip:



Main issue to be addressed (imposed by the technology) : **2 input qu-bits of a non single gate (e.g. CNOT) needs to be adjacent to compute the gate**

- ↪ 2 qu-bits are adjacent if they are either on the same row or on the same column
- ↪ If they are not adjacent, they need to be moved to satisfy this constraint (**routing process**)

## Placement

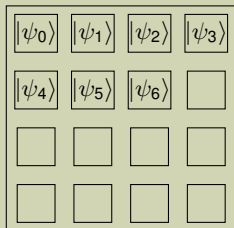
Maps the qu-bits on chip, deriving the initial configuration of the processor

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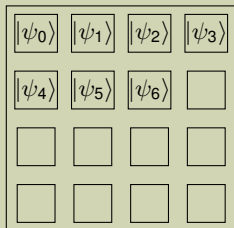


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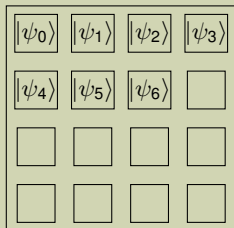
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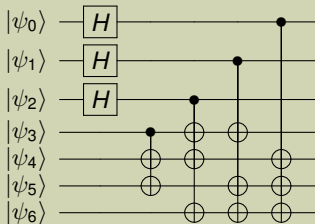
- **Target:** finding the placement minimizing the sum of Manhattan distances over all pairs involved in multiple bits gates
- One possible approach: **Quantum Interaction Graph (QIG)**



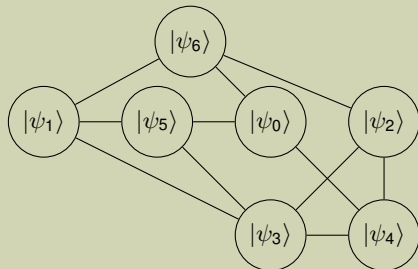
## Quantum Interaction Graph

- QIG purpose: represent the relationships between qu-bits involved in multiple bits gates
- The corresponding symmetric matrix can be used to define a linear programming problem
- The solution of this problem provide a good placement.

Example Circuit:



Quantum Interaction Graph:



Gates can theoretically be all executed simultaneously, but:

## Scheduling Issues

- Data dependencies
- Privileged Writings on each qu-bit
- Out of order execution must preserve the correctness of the computation

# Quantum Compilation: Scheduling

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- Privileged Writings on each qu-bit
- Out of order execution must preserve the correctness of the computation

## Scheduling Policy

- An As Soon As Possible (ASAP) policy is usually employed
  - An operation is performed as soon as the input data are available
- Mainly due to decoherence time constraints, **As Late As Possible (ALAP) policy** is generally preferable in quantum scenarios
- Try to minimize the time between an operation writing a qu-bit and the next operation reading it → reducing the time interval the quantum state needs to be preserved

## ASAP Policy

**C0:**  $\text{.init } |\psi_0\rangle, |\psi_1\rangle, |\psi_2\rangle, |\psi_3\rangle, |\psi_4\rangle, |\psi_5\rangle, |\psi_6\rangle$

**C1:**  $H(|\psi_0\rangle), H(|\psi_1\rangle), H(|\psi_2\rangle), CNOT(|\psi_3\rangle, |\psi_4\rangle), CNOT(|\psi_3, \psi_5\rangle)$

**C2:**  $CNOT(|\psi_2\rangle, |\psi_3\rangle), CNOT(|\psi_2\rangle, |\psi_4\rangle), CNOT(|\psi_2\rangle, |\psi_6\rangle), CNOT(|\psi_1\rangle, |\psi_5\rangle)$

**C3:**  $CNOT(|\psi_1\rangle, |\psi_3\rangle), CNOT(|\psi_1\rangle, |\psi_6\rangle), CNOT(|\psi_0\rangle, |\psi_4\rangle), CNOT(|\psi_0\rangle, |\psi_5\rangle)$

**C4:**  $CNOT(|\psi_0\rangle, |\psi_6\rangle)$

## ASAP Policy

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**C4:** `CNOT( $|\psi_0\rangle, |\psi_6\rangle$ )`

## ALAP Policy

**C0:** `.init  $|\psi_2\rangle$`

**C1:** `.init  $|\psi_1\rangle, |\psi_3\rangle, |\psi_4\rangle, |\psi_5\rangle, |\psi_6\rangle, H(|\psi_2\rangle)$`

**C2:** `.init  $|\psi_0\rangle, H(|\psi_1\rangle)$ , CNOT( $|\psi_3\rangle, |\psi_4\rangle$ ), CNOT( $|\psi_3, \psi_5\rangle$ ), CNOT( $|\psi_2\rangle, |\psi_6\rangle$ )`

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Postpone initialization as late as possible!

Recall: to perform multi qu-bit gates the qu-bits need to be adjacent

↪ If they are not, we need to move them. How?

↪ Using swap gate!

# Quantum Compilation: Routing

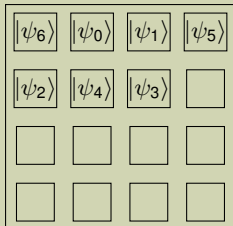
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## Routing Example

Consider this initial placement obtained from our linear programming algorithm:



Now, consider the third cycle using ALAP scheduling policy:

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↪ We need to add a swap between  $|\psi_5\rangle$  and  $|\psi_1\rangle$

↪ .init  $|\psi_0\rangle$ ,  $H(|\psi_1\rangle)$ ,  $CNOT(|\psi_3, |\psi_4\rangle)$ ,  $SWAP(|\psi_1, \psi_5\rangle)$ ,  $CNOT(|\psi_2, |\psi_6\rangle)$

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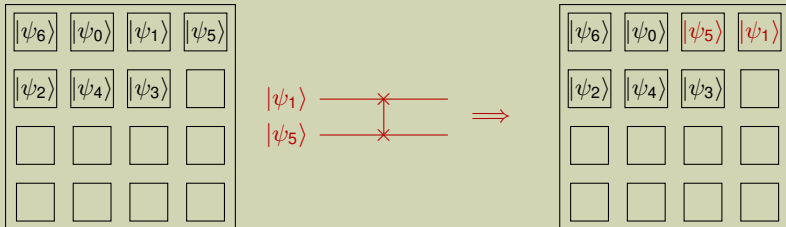
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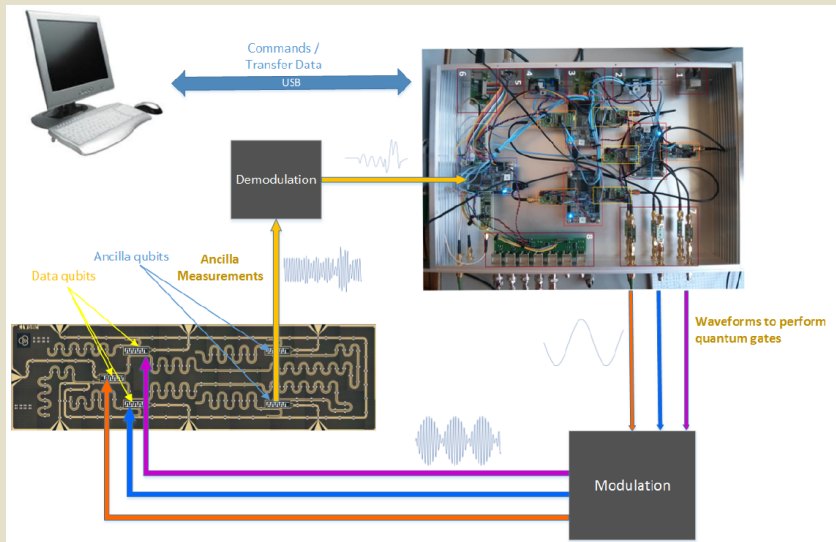
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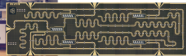
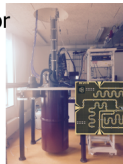
## Technical Challenges

- Qu-bits can preserve their states only at really low temperatures (mK order)
- They need to interact with electronic components → the heat generated by these components should not affect the chip
- Cooling methods:
  1. Heat bath: liquid helium is employed
  2. In 2017, a quantum refrigerator chip based on tunnel effect has been proposed

1<sup>st</sup> floor



Ground floor  
20mK



## Shielding

- The quantum computer needs to be isolated from external electro-magnetic waves
- The quantum processor is shielded with magnetic elements which zero the magnetic field coming from electronic circuitry

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## Wiring

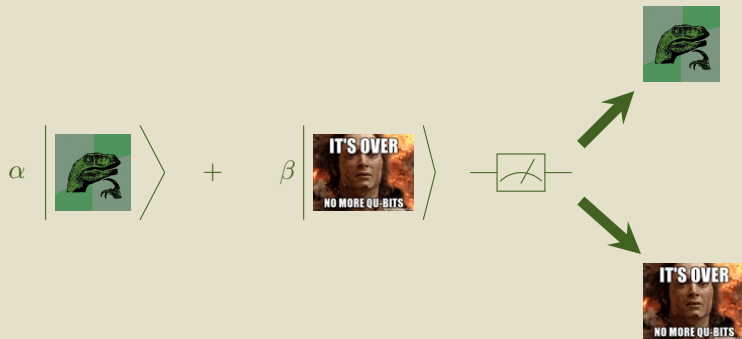
- Wiring: with a lot of qu-bits, placement of wires to perform operations may become complex
- In particular, interferences among different wires is a relevant issue
- Due to:
  1. Wires need to work at extremely low temperatures
  2. Material used to build wires cannot be magnetic



non conventional material is necessary

## Take Home Messages

- Qu-bit superposition state allows to represent simultaneously both 0 and 1
- Qu-bit measurement is probabilistic
- Quantum gates are reversible
- Entanglement: 2 entangled qu-bits are strictly linked, and operations performed on one qu-bit may affect the other one too
- Entanglement can be used for quantum teleportation
- Exponential improvement:  $N$  qu-bits allow to represent  $2^N$  values
- Quantum Algorithm idea: perform computation which yields to high probability of measuring the correct result
- Quantum technology is fragile → error correction is necessary
- Quantum processor are nowadays too limited for practical application
- Quantum computer architecture & compilation
- There are relevant technical challenges to build a quantum computer



- From the previous example, we can see that routing affects the scheduling
- Swaps are additional gates, which introduce new constraints
- But we cannot properly insert swaps if we do not know the scheduling of the operations



**Routing & Scheduling should be performed together given the initial placement**

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## Optimal Solution?

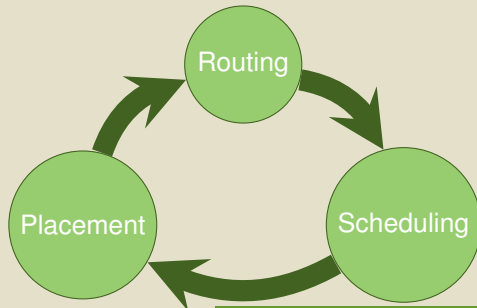
- Routing cost is estimated on the initial placement for all gates
- But the placement of the qu-bits changes during the execution
- The solution **may not be optimal!**
- However, if we consider the temporal dependencies during placement the problem becomes more complex: scheduling of the operations is relevant too!



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Routing & Scheduling should be performed together given the initial placement



- 5 qu-bits IBM processor under the curtains:  
<https://arstechnica.com/science/2016/05/how-ibms-new-five-qubit-universal-quantum-computer-works/>
- IBM Quantum Experience:  
<https://quantumexperience.ng.bluemix.net/qx>
- Quantum Computer Simulator: <http://quantum-studio.net/>

## Generic Qu-Bit Computation

$$|\psi_m\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle, |\psi_a\rangle = |\psi_b\rangle = |0\rangle \Rightarrow |\psi_{in}\rangle = \alpha_0 |000\rangle + \alpha_1 |100\rangle$$

1. Bell pair creation:  $\alpha_0 \frac{1}{\sqrt{2}}(|000\rangle + |011\rangle) + \alpha_1 \frac{1}{\sqrt{2}}(|100\rangle + |111\rangle)$

2. CNOT gate:  $\alpha_0 \frac{1}{\sqrt{2}}(|000\rangle + |011\rangle) + \alpha_1 \frac{1}{\sqrt{2}}(|110\rangle + |101\rangle)$

3. Hadamar gate:

$$\alpha_0 \frac{1}{2}(|000\rangle + |100\rangle + |011\rangle + |111\rangle) + \alpha_1 \frac{1}{2}(|010\rangle - |110\rangle + |001\rangle - |101\rangle)$$

4. Split before measurement:  $\frac{1}{2} |00\rangle \otimes (\alpha_0 |0\rangle + \alpha_1 |1\rangle) + \frac{1}{2} |01\rangle \otimes (\alpha_1 |0\rangle + \alpha_0 |1\rangle) + \frac{1}{2} |10\rangle \otimes (\alpha_0 |0\rangle - \alpha_1 |1\rangle) + \frac{1}{2} |11\rangle \otimes (-\alpha_1 |0\rangle + \alpha_0 |1\rangle)$

5. Decoding:

Measure	$ \psi_b\rangle$	Corrections	Output
00	$\alpha_0  0\rangle + \alpha_1  1\rangle$	No	$\alpha_0  0\rangle + \alpha_1  1\rangle$
01	$\alpha_1  0\rangle + \alpha_0  1\rangle$	X	$\alpha_0  0\rangle + \alpha_1  1\rangle$
10	$\alpha_0  0\rangle - \alpha_1  1\rangle$	Z	$\alpha_0  0\rangle + \alpha_1  1\rangle$
11	$-\alpha_1  0\rangle + \alpha_0  1\rangle$	X,Z	$\alpha_0  0\rangle + \alpha_1  1\rangle$